Applications of Formal Verification

Model Checking: Introduction to SPIN

Dr. Vladimir Klebanov · Dr. Mattias Ulbrich | SS 2015
**SPIN: Previous Lecture vs. This Lecture**

**Previous lecture**
*SPIN* appeared as a PROMELA *simulator*

**This lecture**
Intro to *SPIN* as a *model checker*
What Does A Model Checker Do?

A Model Checker (MC) is designed to prove the user wrong.

MC tries its best to *find a counter example* to the correctness properties. It is tuned for that.

MC does not try to prove correctness properties. It tries the opposite.
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But why then can a MC also *prove* correctness properties?
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MC does not try to prove correctness properties. It tries the opposite.

But why then can a MC also prove correctness properties?

MC’s search for counter examples is exhaustive.

⇒ Finding no counter example proves stated correctness properties.
What does ‘exhaustive search’ mean here?

exhaustive search  
=  
resolving non-determinism in all possible ways
What does ‘exhaustive search’ mean here?

exhaustive search

= resolving non-determinism in all possible ways

For model checking PROMELA code, two kinds of non-determinism to be resolved:

- **explicit, local:** if/do statements
  
  :: guardX -> ....
  
  :: guardY -> ....

- **implicit, global:** scheduling of concurrent processes
  
  (see next lecture)
Model Checker for This Course: **SPIN**

**SPIN**: “Simple Promela Interpreter”
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If this was all, you would have seen most of it already. The name is a serious understatement!
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Main functionality of **SPIN**:  
- simulating a model (randomly/interactively)  
- generating a verifier
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**verifier** generated by **SPIN** is a C program performing...
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Main functionality of SPIN:
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- generating a \textit{verifier}

\textit{Verifier} generated by SPIN is a C program performing \textit{model checking}:
- exhaustively checks PROMELA model against correctness properties
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\textit{Verifier} generated by \textsf{SPIN} is a C program performing \textit{model checking}:
- exhaustively \textit{checks} \texttt{PROMELA model} against correctness properties
- in case the check is negative: generates a \textit{failing run} of the model
Model Checker for This Course: SPIN

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Main functionality of SPIN:
- simulating a model (randomly/interactively/guided)
- generating a verifier

Verifier generated by SPIN is a C program performing model checking:
- exhaustively checks PROMELA model against correctness properties
- in case the check is negative: generates a failing run of the model, to be simulated by SPIN
**SPIN Workflow: Overview**

1. **Model**
   - `name.pml`

2. **Correctness Properties**

3. **Spin**
   - `-a`
   - `-i`
   - `-t`

4. **Verifier**
   - `verifier`
   - `pan.c`

5. **C Compiler**
   - `C compiler`

6. **Executable Verifier**
   - `pan`

7. **Run**
   - `name.pml.trail`
   - `failing`
   - `interactive / guided simulation`

8. **Results**
   - `errors: 0`
   - `either`
   - `0`

9. **Output**
   - `SS 2015 6/34`
Plain Simulation with SPIN

model
name.pml

correctness
properties

SPIN

verifier
pan.c

C
compiler

executable
verifier
pan

-f
random/
interactive/
guided
simulation

failing
run
name.pml.trail

"errors: 0"
Rehearsal: Simulation Demo

- run example, random and interactive
  
  `interleave.pml, zero.pml`
Model Checking with SPIN

- model name.pml
- correctness properties

SPIN

- verifier 
- pan.c

C compiler

executable verifier

pan

random/ interactive/ guided simulation

verifier pan.c

either

failing run name.pml.trail

"errors: 0"

Meaning of Correctness wrt. Properties

Given PROMELA model $M$, and correctness properties $C_1, \ldots, C_n$.

- Be $R_M$ the set of all possible runs of $M$.
- For each correctness property $C_i$, $R_{M,C_i}$ is the set of all runs of $M$ satisfying $C_i$. ($R_{M,C_i} \subseteq R_M$)
- $M$ is correct wrt. $C_1, \ldots, C_n$ iff $(R_{M,C_1} \cap \ldots \cap R_{M,C_n}) = R_M$.
- If $M$ is not correct, then each $r \in (R_M \setminus (R_{M,C_1} \cap \ldots \cap R_{M,C_n}))$ is a counter example.
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We know how to write models $M$. 
Meaning of Correctness wrt. Properties

Given PROMELA model $M$, and correctness properties $C_1, \ldots, C_n$.

- Be $R_M$ the set of all possible runs of $M$.
- For each correctness property $C_i$, $R_{M,C_i}$ is the set of all runs of $M$ satisfying $C_i$.
  \[(R_{M,C_1} \subseteq R_M)\]
- $M$ is correct wrt. $C_1, \ldots, C_n$ iff \((R_{M,C_1} \cap \ldots \cap R_{M,C_n}) = R_M\).
- If $M$ is not correct, then each $r \in (R_M \setminus (R_{M,C_1} \cap \ldots \cap R_{M,C_n}))$ is a counter example.

We know how to write models $M$.
But how to write Correctness Properties?
Correctness properties can be stated (syntactically) within or outside the model. Stating properties within the model, using assertion statements, meta labels, and end labels. Stating properties outside the model, using never claims or temporal logic formulas.
Correctness properties can be stated (syntactically) within or outside the model.
Stating Correctness Properties

Correctness properties can be stated (syntactically) within or outside the model.

- stating properties within the model, using assertion statements
Stating Correctness Properties

Correctness properties can be stated (syntactically) **within** or **outside** the model.

- **Stating properties within the model**, using:
  - assertion statements
  - meta labels
    - `end` labels
    - `accept` labels
    - `progress` labels

**model** name.pml

**correctness properties**
Correctness properties can be stated (syntactically) **within** or **outside** the model.

**stating properties within the model, using**
- assertion statements
- meta labels
  - end labels
  - accept labels
  - progress labels

**stating properties outside the model, using**
- never claims
- temporal logic formulas
Correctness properties can be stated (syntactically) within or outside the model.

stating properties within the model, using
  - *assertion statements* (today)
  - meta labels
    - *end labels* (today)
    - accept labels
    - progress labels

stating properties outside the model, using
  - never claims
  - temporal logic formulas
Assertion Statements

**Definition (Assertion Statements)**

Assertion statements in PROMELA are statements of the form

```
assert (expr)
```

were `expr` is any PROMELA expression.
Assertion Statements

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\[ \text{assert} (\text{expr}) \]

were \( \text{expr} \) is any PROMELA expression.

Typically, \( \text{expr} \) is of type \texttt{bool}.
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Assertion statements can appear anywhere where a PROMELA statement is expected.
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```
assert (expr)
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were `expr` is any PROMELA expression.

Typically, `expr` is of type `bool`.

Assertion statements can appear anywhere where a PROMELA statement is expected.

```c
... stmt1;
assert (max == a);
stmt2;
...```

**Assertion Statements**

**Definition (Assertion Statements)**

Assertion statements in PROMELA are statements of the form

\[
\text{assert}(\text{expr})
\]

were \textit{expr} is any PROMELA expression.

Typically, \textit{expr} is of type \texttt{bool}.

Assertion statements can appear anywhere where a PROMELA statement is expected.

\[
\ldots
\text{stmt1;}
\text{assert}(\text{max} == \text{a});
\text{stmt2;}
\ldots
\]

\[
\ldots
\text{if}
:: \text{b1} \rightarrow \text{stmt3; assert}(\text{x} < \text{y})
:: \text{b2} \rightarrow \text{stmt4}
\ldots
\]
Meaning of **Boolean Assertion Statements**

assert(*expr*)

- has no effect if *expr* evaluates to true
- triggers an error message if *expr* evaluates to false

This holds in both, simulation and model checking mode.
Meaning of **General Assertion Statements**

```plaintext
assert(expr)
```

- has no effect if `expr` evaluates to non-zero value
- triggers an error message if `expr` evaluates to 0

This holds in both, simulation and model checking mode.
Meaning of General Assertion Statements

assert\((expr)\)
- has no effect if \(expr\) evaluates to non-zero value
- triggers an error message if \(expr\) evaluates to 0

This holds in both, simulation and model checking mode.

Recall:

\textbf{bool true false} is syntactic sugar for
assert \((expr)\)
- has no effect if \(expr\) evaluates to non-zero value
- triggers an error message if \(expr\) evaluates to 0

This holds in both, simulation and model checking mode.

Recall:

```
bool true false   is syntactic sugar for
bit   1   0
```
Meaning of **General Assertion** Statements

```markdown
assert(expr)
- has no effect if `expr` evaluates to **non-zero value**
- triggers an error message if `expr` evaluates to 0
```

This holds in both, simulation and model checking mode.

Recall:

```markdown
bool true false  is syntactic sugar for
bit 1 0
⇒ general case covers Boolean case
```
Instead of using ‘printf’ s for Debugging ...

/* after choosing a, b from {1, 2, 3} */
if
  :: a >= b -> max = a;
  :: a <= b -> max = b;
fi;
printf("the maximum of %d and %d is %d\n", a, b, max);
Instead of using ‘printf’s for Debugging ...

/* after choosing a,b from {1,2,3} */
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  :: a >= b -> max = a;
  :: a <= b -> max = b;
fi;
printf("the maximum of %d and %d is %d\n", a, b, max);

Command Line Execution
(simulate, inject faults, add assertion, simulate again)

> spin max.pml
... we can employ **Assertions**

quoting from file `max.pml`:

```plaintext
/* after choosing a,b from \{1,2,3\} */
if
  :: a >= b -> max = a;
  :: a <= b -> max = b;
fi;
assert( a > b -> max == a : max == b )
```
... we can employ **Assertions**

quoting from file `max.pml`:

```pml
/* after choosing a,b from {1,2,3} */
if
    :: a >= b -> max = a;
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fi;
assert( a > b -> max == a : max == b )
```

Now, we have a first example with a formulated **correctness property**.
... we can employ **Assertions**

quoting from file `max.pml`:

```pml
/* after choosing a,b from {1,2,3} */
if
  :: a >= b -> max = a;
  :: a <= b -> max = b;
fi;
assert( a > b -> max == a : max == b )
```

Now, we have a first example with a formulated **correctness property**.

We can do **model checking**, for the first time!
Generate Verifier in C

Command Line Execution

Generate Verifier in C
> spin -a max.pml

SPIN generates Verifier in C, called pan.c
(plus helper files)
Compile To Executable Verifier

Command Line Execution

```
> gcc -o pan pan.c
```

C compiler generates executable verifier `pan` historically "protocol analyzer", now "process analyzer"
Compile To Executable Verifier

Command Line Execution

compile to executable verifier

> gcc -o pan pan.c

C compiler generates executable verifier pan
Compile To Executable Verifier

Command Line Execution

*compile to executable verifier*

> gcc -o pan pan.c

C compiler generates executable verifier **pan**

**pan**: historically “protocol analyzer”, now “process analyzer”
Run Verifier (= Model Check)

Executable verifier **pan**

either

"errors: 0"

or

failing run max.pml.trail

Command Line Execution

```
run verifier pan
> ./pan
```
Run Verifier (= Model Check)

Command Line Execution

```
run verifier pan
> ./pan
```

- prints "errors: 0"
Run Verifier (= Model Check)

executable verifier pan

either

“errors: 0”
or

failing run max.pml.trail

Command Line Execution

run verifier pan

> ./pan

- prints “errors: 0” \(\Rightarrow\) Correctness Property verified!
Run Verifier (= Model Check)

Executable verifier `pan`

either

"errors: 0"

or

failing run `max.pml.trail`

Command Line Execution

```
run verifier pan
> ./pan
```

- prints "errors: 0", or
- prints "errors: n" (n > 0)
Run Verifier (= Model Check)

**Command Line Execution**

```
run verifier pan
> ./pan
```

- prints "errors: 0", or
- prints "errors: n" (n > 0)  ⇒ counter example found!
Run Verifier (= Model Check)

executable verifier pan

either

“errors: 0”

or

failing run max.pml.trail

Command Line Execution

run verifier pan

> ./pan

- prints “errors: 0”, or
- prints “errors: n” \( (n > 0) \) ⇒ counter example found!
records failing run in max.pml.trail
Guided Simulation

To examine failing run: employ simulation mode, “guided” by trail file.

Command Line Execution

inject a fault, re-run verification, and then:

> spin -t -p -l max.pml
can look like:

Starting P with pid 0

1: proc 0 (P) line 8 "max.pml" (state 1) [a = 1]
   P(0):a = 1

2: proc 0 (P) line 14 "max.pml" (state 7) [b = 2]
   P(0):b = 2

3: proc 0 (P) line 23 "max.pml" (state 13) [((a<=b))]
3: proc 0 (P) line 23 "max.pml" (state 14) [max = a]
   P(0):max = 1

spin: line 25 "max.pml", Error: assertion violated
spin: text of failed assertion:
   assert(( ((a>b)) -> ((max==a)) : ((max==b)) ))
can look like:

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assignments in the run
can look like:

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spin: line 25 "max.pml", Error: assertion violated
spin: text of failed assertion:
   assert(( ((a>b)) -> ((max==a)) : ((max==b)) ))

assignments in the run
values of variables whenever updated
What did we do so far?

following whole cycle (most primitive example, assertions only)

- model `name.pml`
- correctness properties

**SPIN**

- `-a`
- `-i`
- `-t`

**verifier**

- `pan.c`

**C compiler**

**executable verifier**

- `pan`

**failing run**

- `name.pml.trail`

**random/ interactive / guided simulation**

```
errors: 0
```

- either
- or

"errors: 0"
What did we do so far?

following whole cycle (most primitive example, assertions only)

model

name.pml
correctness

properties

SPIN

verifier
pan.c

C
compiler

executable
verifier
pan

verifier

pan.c

failing
run
name.pml.trail

"errors: 0"
Further Examples: Integer Division

```c
int dividend = 15;
int divisor = 4;
int quotient, remainder;

quotient = 0;
remainder = dividend;
do
  :: remainder > divisor ->
    quotient++;
    remainder = remainder - divisor
  :: else ->
    break
od;
printf("%d divided by %d = %d, remainder = %d\n", dividend, divisor, quotient, remainder);
```
Further Examples: Integer Division

```c
int dividend = 15;
int divisor = 4;
int quotient, remainder;

quotient = 0;
remainder = dividend;

while (remainder > divisor) {
    quotient++;
    remainder = remainder - divisor;
}

printf("%d divided by %d = %d, remainder = %d\n", dividend, divisor, quotient, remainder);
```

simulate, put assertions, verify, change values, ...
Further Examples: Greatest Common Divisor

```c
int x = 15, y = 20;
int a, b;
a = x; b = y;
do
    :: a > b -> a = a - b
    :: b > a -> b = b - a
    :: a == b -> break
od;
printf("The GCD of %d and %d = %d\n", x, y, a)
```
Further Examples: Greatest Common Divisor

```c
int x = 15, y = 20;
int a, b;
a = x; b = y;
do
  :: a > b -> a = a - b
  :: b > a -> b = b - a
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full functional verification not possible here (why?)
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full functional verification not possible here (why?)
still, assertions can perform sanity check
Further Examples: Greatest Common Divisor

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int x = 15, y = 20;
int a, b;
a = x; b = y;
do
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    :: a == b -> break
od;
printf("The GCD of %d and %d = %d\n", x, y, a)
```

full functional verification not possible here (why?)

still, assertions can perform **sanity check**

⇒ typical for model checking
Typical Command Lines

typical command line sequences:
random simulation
   spin name.pml
Typical Command Lines

typical command line sequences:

random simulation
  spin name.pml

interactive simulation
  spin -i name.pml
Typical Command Lines

typical command line sequences:
random simulation
   spin name.pml
interactive simulation
   spin -i name.pml
model checking
   spin -a name.pml
   gcc -o pan pan.c
   ./pan
Typical Command Lines

typical command line sequences:
random simulation
   spin name.pml
interactive simulation
   spin -i name.pml
model checking
   spin -a name.pml
gcc -o pan pan.c
   ./pan
and in case of error
   spin -t -p -l -g name.pml
Spin Reference Card

Ben-Ari produced Spin Reference Card, summarizing

- typical command line sequences
- options for
  - SPIN
  - gcc
  - pan
- PROMELA
  - datatypes
  - operators
  - statements
  - guarded commands
  - processes
  - channels
- temporal logic syntax
Why SPIN?

- SPIN targets software, instead of hardware verification
- based on standard theory of $\omega$-automata and linear temporal logic
- 2001 ACM Software Systems Award (other winning software systems include: Unix, TCP/IP, WWW, Tcl/Tk, Java)
- used for safety critical applications
- distributed freely as research tool, well-documented, actively maintained, large user-base in academia and in industry
- annual SPIN user workshops series held since 1995
Why SPIN? (Cont’d)

- PROMELA and SPIN are rather simple to use
- good to understand a few system really well, rather than many systems poorly
- availability of good course book (Ben-Ari)
- availability of front end JSPIN (also Ben-Ari)
What is JSPIN?

- graphical user interface for SPIN
- developed for pedagogical purposes
- written in Java
- simple user interface
- SPIN options automatically supplied
- fully configurable
- supports graphics output of transition system
What is JSPI\textsc{N}?

- graphical user interface for SP\textsc{IN}
- developed for pedagogical purposes
- written in Java
- simple user interface
- SP\textsc{IN} options automatically supplied
- fully configurable
- supports graphics output of transition system
- makes back-end calls transparent
**Command Line Execution**

*calling JSPIN*

> java -jar /usr/local/jSpin/jSpin.jar

*(with path adjusted to your setting)*

*or use shell script:*

> jspin
**JSPIN Demo**

---

**Command Line Execution**

*calling JSPIN*

```bash
> java -jar /usr/local/jSpin/jSpin.jar
```

*(with path adjusted to your setting)*

*or use shell script:*

```bash
> jspin
```

play around with similar examples ...
Catching A Different Type of Error

quoting from file **max2.pml**:

```pml
/* after choosing a, b from {1, 2, 3} */
if
  :: a >= b -> max = a;
  :: b <= a -> max = b;
fi;
printf("the maximum of %d and %d is %d\n",
  a, b, max);
```

simulate a few times ⇒ crazy "timeout" message sometimes
generate and execute pan ⇒ reports "errors: 1"
Catching A Different Type of Error

quoting from file max2.pml:

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Catching A Different Type of Error

quoting from file max2.pml:

```c
/* after choosing a,b from {1,2,3} */
if
  :: a >= b -> max = a;
  :: b <= a -> max = b;
fi;
printf("the maximum of %d and %d is %d\n", 
       a, b, max);
```

simulate a few times
⇒ crazy “timeout” message sometimes

generate and execute pan
⇒ reports “errors: 1”

????
Catching A Different Type of Error

quoting from file max2.pml:

```c
/* after choosing a,b from {1,2,3} */
if
    :: a >= b -> max = a;
    :: b <= a -> max = b;
fi;
printf("the maximum of %d and %d is %d\n", 
a, b, max);
```

simulate a few times
⇒ crazy “timeout” message sometimes

generate and execute pan
⇒ reports “errors: 1”

Note: no assert in max2.pml.
Further inspection of pan output:

...  
pan: invalid end state (at depth 1)  
pan: wrote max2.pml.trail  
...
Legal and Illegal Blocking

A process may legally block, **as long as some other process can proceed.**
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in max1.pml, no process can take over.
Valid End States

Definition (Valid End State)
An end state of a run is valid iff the location counter of each processes is at an end location.
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- each location marked with an end label: “endxxx:”

End labels are not useful in max1.pml, but elsewhere, they are. Example: end.pml
Literature for this Lecture

Ben-Ari  Chapter 2, Sections 4.7.1, 4.7.2