
Formal Specification of Software

Steam Boiler Control

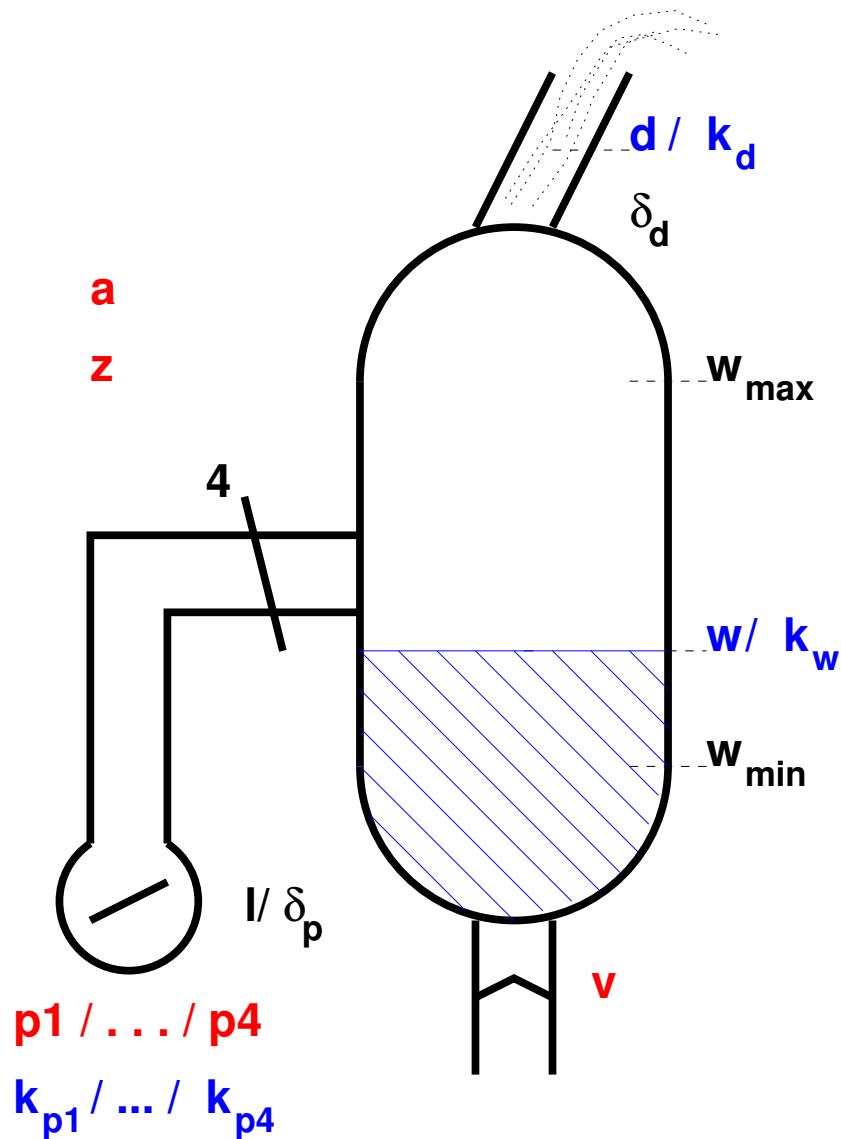
An Example in ASM Formalisation

Bernhard Beckert



UNIVERSITÄT KOBLENZ-LANDAU

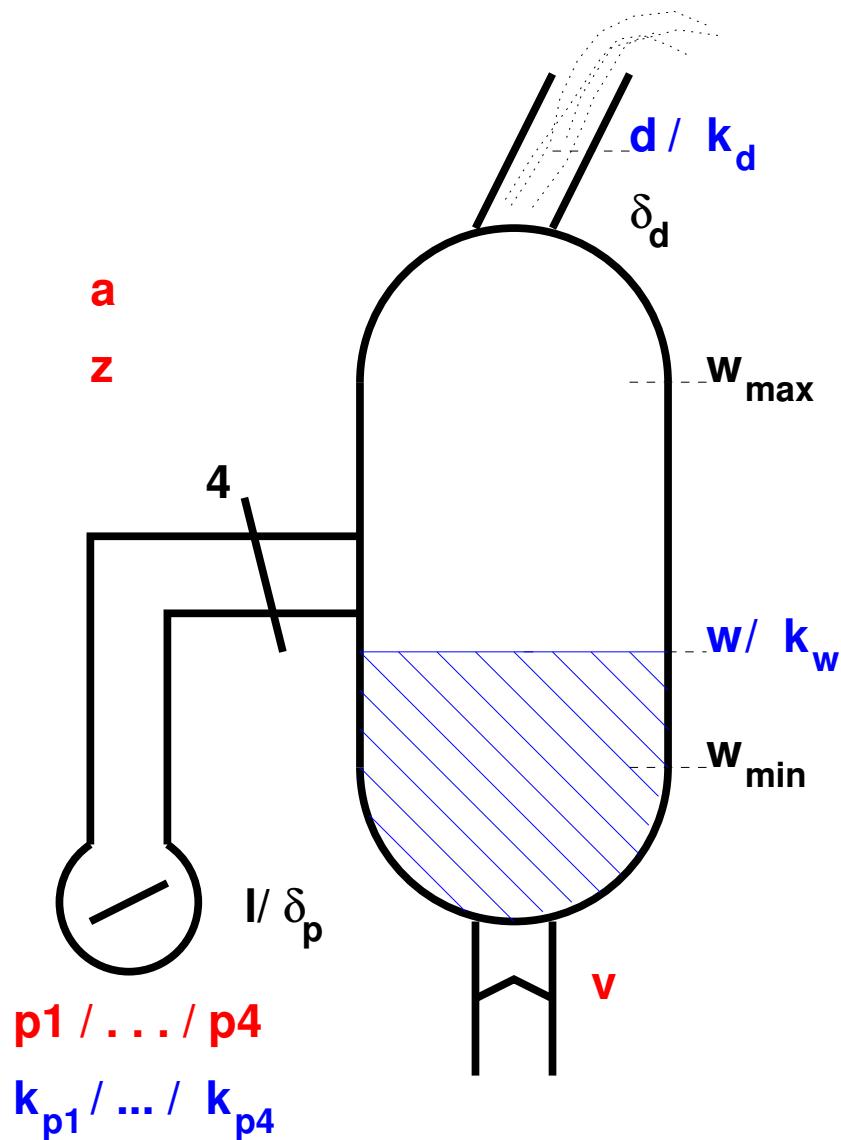
Steam Boiler Control: Scenario



System Components

- steam boiler
- water level measuring device
- four pumps
- four pump controllers
- steam quantity measuring device
- valve for emptying the boiler

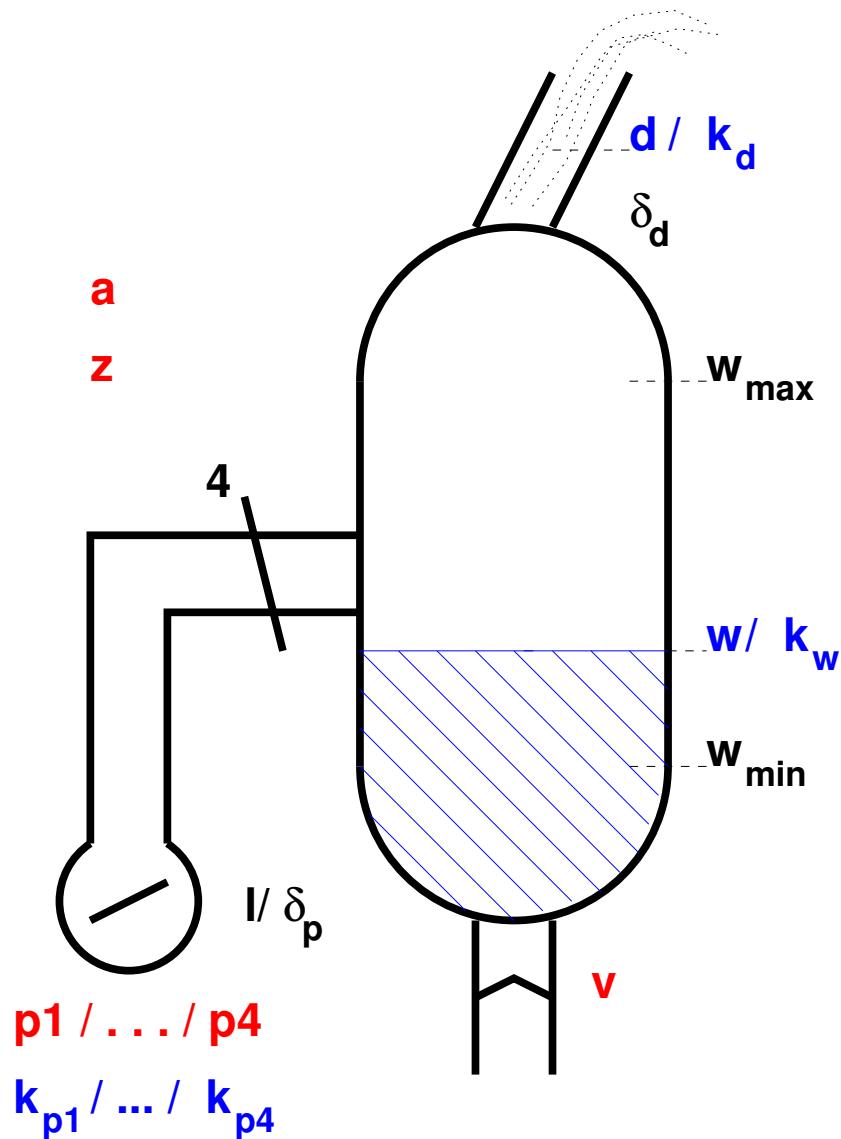
Steam Boiler Control: Scenario



Physical constants

w_{min}	minimal water level
w_{max}	maximal water level
l	water amount per pump
d_{max}	maximal quantity of steam exiting the boiler
δ_p	error in the value of l
δ_d	error in steam measurement

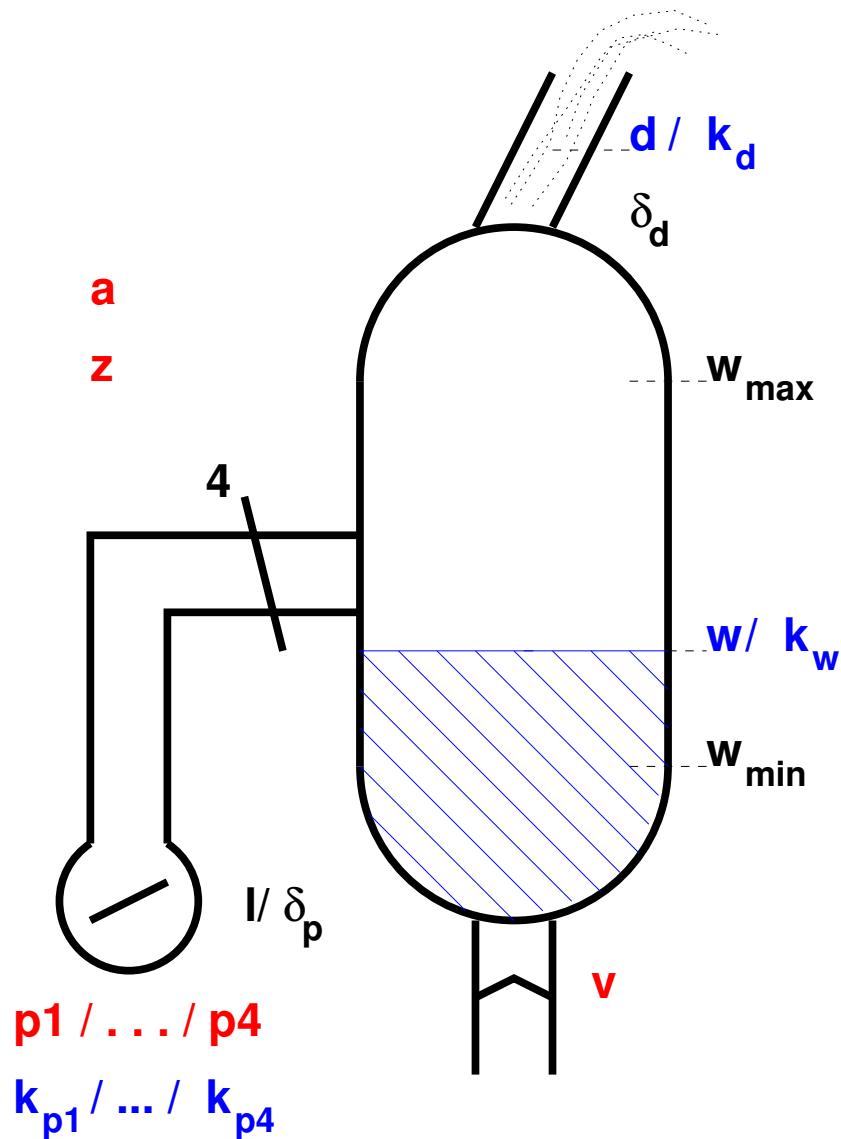
Steam Boiler Control: Scenario



Measured values

w	water level
d	amount of steam exiting the boiler
$k_p(i)$	pump i works/broken
k_w	water level measuring device works/broken
k_d	steam amount measuring device works/broken

Steam Boiler Control: Scenario



Control values

$p(i)$ pump i on/off

v valve open/closed

a boiler on/off

z state

init/norm/broken/stop

Steam Boiler with ASMs

Restrictions

- Real-time aspects not modelled
- Communication between devices not modelled

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- Real-time aspects not modelled
- Communication between devices not modelled

Measured values

Modelled as functions that are changed externally

Control values

Modelled as functions that are read externally

Steam Boiler with ASMs: Two Versions

First version

The possibility that devices are broken is not modelled

States: *init, normal, stop*

Second version

The possibility that devices are broken is included in the model

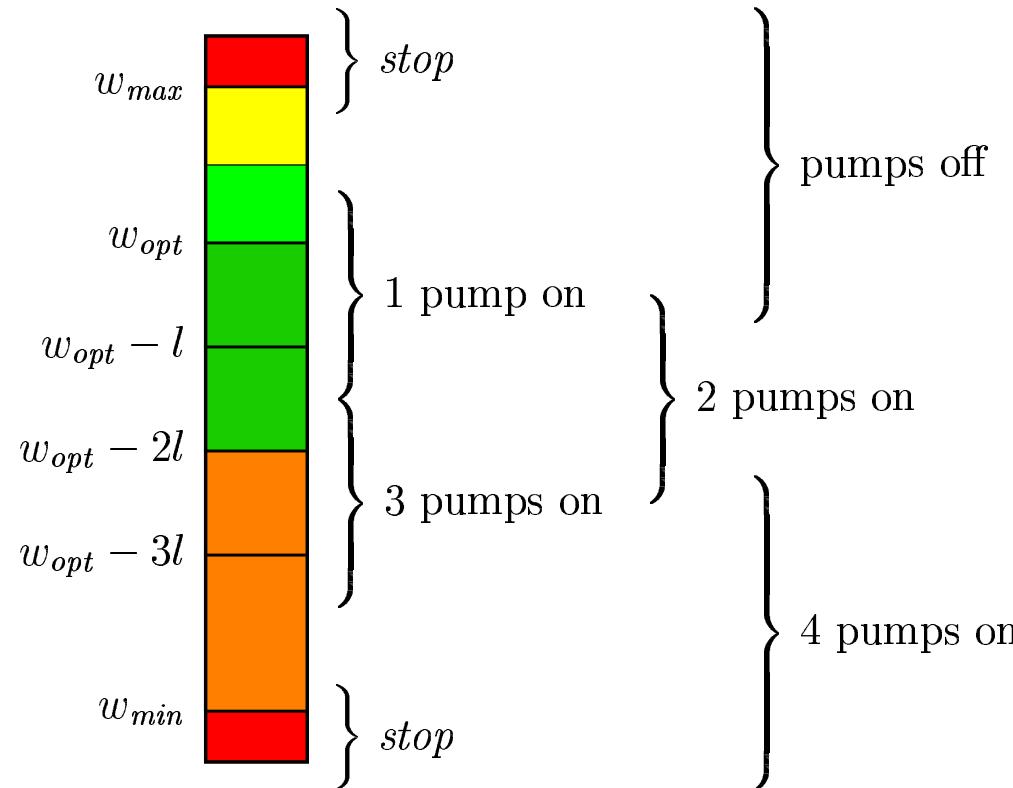
Additional state: *broken*

First Version: Strategy for Filling

Additional constant w_{opt}

Optimal water level

Strategy



First Version: Vocabulary

Universes

$$state = \{init, norm, stop\}$$

$$openClosed = \{open, closed\}$$

$$water = \mathbb{N}$$

$$pumps = \{1, 2, 3, 4\}$$

$$onOff = \{on, off\}$$

First Version: Vocabulary

Universes

state = {*init, norm, stop*}

openClosed = {*open, closed*}

water = \mathbb{N}

pumps = {1,2,3,4}

onOff = {*on, off*}

Note

These are unary boolean functions; they define a type/class

First Version: Vocabulary

Dynamic functions

$p :$	pumps → onOff	controlling the pumps
$v :$	→ openClosed	controlling the steam
$a :$	→ onOff	valve
$z :$	→ state	controlling the boiler
		boiler state

External functions

$w :$	→ water	water level
$d :$	→ water	steam exiting boiler

Static functions

$+, -, * :$	$\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$	arithmetic
$<, \leq :$	$\mathbb{N} \times \mathbb{N} \rightarrow \text{Boole}$	ordering
$w_{max}, w_{min}, w_{opt}, l, d_{max}$	$\rightarrow \mathbb{N}$	physical constants

Initial State

$a = off$

$z = init$

Rule Initialisation

```
if  $\neg(z = init)$  then
    skip
else
    if  $0 < d$  then
         $z := stop$ 
    else if  $w < w_{min} + d_{max}$  then
        par
             $v := closed$ 
             $p(i) := on \quad (i = 1..4)$ 
        endpar
    else if  $w_{max} < w$  then
        par
             $v := open$ 
             $p(i) := off \quad (i = 1..4)$ 
        endpar
```

```
else
    par
         $z := norm$ 
         $v := closed$ 
         $a := on$ 
         $p(i) := off \quad (i = 1..4)$ 
    endpar
endif endif endif
endif
```

Rule Normal

```
if  $\neg(z = norm)$  then
    skip
else
    if  $w_{max} < w \vee w < w_{min}$  then
        par
            a := off
            z := stop
        endpar
    else
        par
            if  $w \leq w_{opt}$  then  $p(1) := on$  else  $p(1) := off$  endif
            if  $w \leq w_{opt} - l$  then  $p(2) := on$  else  $p(2) := off$  endif
            if  $w \leq w_{opt} - (2 * l)$  then  $p(3) := on$  else  $p(3) := off$  endif
            if  $w \leq w_{opt} - (3 * l)$  then  $p(4) := on$  else  $p(4) := off$  endif
        endpar
    endif
endif
```

Rule Control

par

Initialisation

Normal

endpar

Second Version: Vocabulary

Universes

state = {*init*, *norm*, *broken*, *stop*}

openClosed = {*open*, *closed*}

water = \mathbb{N}

pumps = {1, 2, 3, 4}

onOff = {*on*, *off*}

worksBroken = {*works*, *broken*}

Second Version: Vocabulary

Dynamic functions

p :	pumps	\rightarrow onOff	controlling the pumps
v :		\rightarrow openClosed	controlling steam valve
a :		\rightarrow onOff	controlling the boiler
z :		\rightarrow state	boiler state
s_{min}, s_{max} :		\rightarrow water	estimated water level
n_p :		\rightarrow pumps	number of active pumps

External functions

w :	\rightarrow water	water level
d :	\rightarrow water	steam exiting boiler
k_p :	pumps \rightarrow worksBroken	pump works/broken
k_w :	\rightarrow worksBroken	water level device
k_d :	\rightarrow worksBroken	steam amount device

Second Version: Vocabulary

Static functions

$+, -, *, \min :$	$\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$	arithmetic
$<, \leq :$	$\mathbb{N} \times \mathbb{N} \rightarrow \text{Boole}$	ordering
$w_{max}, w_{min}, l :$	$\rightarrow \mathbb{N}$	physical constants
$d_{max}, \delta_p, \delta_d :$	$\rightarrow \mathbb{N}$	physical constants
<i>optPumps</i> :	water \times water \rightarrow pumps	optimal pump number
<i>numWorking</i> :	$\mathbb{N} \times \text{worksBroken}^4 \rightarrow \mathbb{N}$	number of working pumps
<i>controlPumps</i> : pumps ² \times worksBroken ⁴ \rightarrow onOff		control for each pump

Second Version: Vocabulary

Static function $optPumps$ **(encodes the strategy)**

$optPumps(w_1, w_2)$ = **optimal number of pumps for
water level between w_1 and w_2**

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Static function $numWorking$

$numWorking(i, k_1, k_2, k_3, k_4)$ = $\#\{j \mid j \leq i \wedge k_j = works\}$

Second Version: Vocabulary

Static function $optPumps$ **(encodes the strategy)**

$optPumps(w_1, w_2)$ = **optimal number of pumps for water level between w_1 and w_2**

Static function $numWorking$

$numWorking(i, k_1, k_2, k_3, k_4) = \#\{j \mid j \leq i \wedge k_j = works\}$

Static function $controlPumps$

$controlPumps(i, n_{opt}, k_1, k_2, k_3, k_4) =$
$$\begin{cases} on & \text{if } numWorking(i - 1, k_1, k_2, k_3, k_4) < n_{opt} \\ off & \text{otherwise} \end{cases}$$

Rule Initialisation

```
if  $\neg(z = init)$  then
    skip
else
    if  $0 < d \vee k_w = broken$ 
         $\vee k_d = broken$  then
             $z := stop$ 
    else if  $w < w_{min} + d_{max}$  then
        par
             $v := closed$ 
             $p(i) := on \quad (i = 1..4)$ 
        endpar
    else if  $w_{max} < w$  then
        par
             $v := open$ 
             $p(i) := off \quad (i = 1..4)$ 
        endpar
```

```
else
    par
         $z := norm$ 
         $v := closed$ 
         $s_{min} := w$ 
         $s_{max} := w$ 
         $n_p := 0$ 
         $p(i) := off \quad (i = 1..4)$ 
    endpar
endif endif endif
endif
```

Rule *NormBroken*

```
if  $\neg(z = norm \vee z = broken)$  then
  skip
else
  if  $k_w = works$  then
    let  $min = w, max = w, z_{val} = norm$  in ControlPumps endlet
  else if  $k_d = works$  then
    let  $min = s_{min} - d + n_p \cdot l - \delta_d - n_p \cdot \delta_p,$ 
         $max = s_{max} - d + n_p \cdot l + \delta_d + n_p \cdot \delta_p,$ 
         $z_{val} = broken$ 
    in ControlPumps endlet
  else
    par
      z := stop
      a := off
    endpar
  endif endif
endif
```

Rule *ControlPumps*

```
if  $min < w_{min} \vee w_{max} < max$  then
  par
    z := stop
    a := off
  endpar
else
  let  $n_{opt} = optPumps(min, max)$  in
    par
      p(i) := controlPumps( $i, n_{opt}, k_p(1), \dots, k_p(4)$ ) ( $i = 1..4$ )
       $n_p := \min(n_{opt}, numWorking(4, k_p(1), \dots, k_p(4)))$ 
       $s_{min} := min$ 
       $s_{max} := max$ 
      z :=  $z_{val}$ 
    endpar
  endlet
endif
```

Rule Control

```
par  
  Initialisation  
  NormBroken  
endpar
```

Alternative Solution: Vocabulary

Universes

state = $\{init, norm, broken, stop\}$

openClosed = $\{open, closed\}$

water = \mathbb{N}

pumps = $\{1, 2, 3, 4\}$

onOff = $\{on, off\}$

worksBroken = $\{works, broken\}$

waitCompute = $\{wait, compute\}$

Alternative Solution: Vocabulary

Additional dynamic functions

$i : \rightarrow$	pumps	current pump
$f : \rightarrow$	waitCompute	next cycle

Meaning of function f

$f = compute$: **Control the pumps**

$f = wait$: **Measurement**

Alternative: Rule *Initialisation*

```
if  $\neg(z = init)$  then
  skip
else
  if  $0 < d \vee k_w = broken$ 
     $\vee k_d = broken$  then
       $z := stop$ 
  else if  $w < w_{min} + d_{max}$  then
    par
       $v := closed$ 
       $p(i) := on \quad (i = 1..4)$ 
       $f := wait$ 
    endpar
  else if  $w_{max} < w$  then
```

```
par
   $v := open$ 
   $p(i) := off \quad (i = 1..4)$ 
   $f := wait$ 
endpar
else
  par
     $z := norm$ 
     $f := wait$ 
     $v := closed$ 
     $s_{min} := w$ 
     $s_{max} := w$ 
     $n_p := 0$ 
     $p(i) := off \quad (i = 1..4)$ 
  endpar
endif endif endif endif
```

Alternative: Rule *NormBroken* (1)

```
if  $\neg((z = norm \vee z = broken) \wedge f = wait)$  then
  skip
else
  if  $k_w = works$  then
    par
       $s_{min} := w$ 
       $s_{max} := w$ 
       $z := norm$ 
       $f := compute$ 
       $i := 1$ 
       $n_p := 0$ 
    endpar
```

Alternative: Rule *NormBroken* (2)

```
else if  $k_d = works$  then
    par
         $s_{min} := s_{min} - d + n_p \cdot l - \delta_d - n_p \cdot \delta_p$ 
         $s_{max} := s_{max} - d + n_p \cdot l + \delta_d + n_p \cdot \delta_p$ 
         $z := broken$ 
         $f := compute$ 
         $i := 1$ 
         $n_p := 0$ 
    endpar
else
    par
         $z := stop$ 
         $a := off$ 
    endpar
endif endif
endif
```

Alternative: Rule *ControlPumps* (1)

```
if  $\neg((z = norm \vee z = broken) \wedge f = compute)$  then
  skip
else
  if  $s_{min} < w_{min} \vee w_{max} < s_{max}$  then
    par
      z := stop
      a := off
    endpar
```

Alternative: Rule *ControlPumps* (2)

```
else
  par
    if  $n_p < optPumps(s_{min}, s_{max}) \wedge k_p(i) = works$  then
      par
         $p(i) := on$ 
         $n_p := n_p + 1$ 
      endpar
    else
       $p(i) := off$ 
    endif
    if  $i < 4$  then
       $i := i + 1$ 
    else
       $f := wait$ 
    endif
  endpar
endif
```

Alternative: Rule Control

```
par
  Initialisation
  NormBroken
  ControlPumps
endpar
```