
Introduction to Artificial Intelligence

First-order Logic

(Logic, Deduction, Knowledge Representation)

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Outline

- **Why first-order logic?**
- **Syntax and semantics of first-order logic**
- **Fun with sentences**
- **Wumpus world in first-order logic**

Pros and Cons of Propositional Logic

- 😊 Propositional logic is **declarative**:
pieces of syntax correspond to facts
- 😊 Propositional logic allows partial / disjunctive / negated information
(unlike most data structures and databases)
- 😊 Propositional logic is **compositional**:
meaning of $B_{1,1} \wedge P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$
- 😊 Meaning in propositional logic is **context-independent**
(unlike natural language, where meaning depends on context)
- 😞 Propositional logic has very limited expressive power
(unlike natural language)

Example:

Cannot say “pits cause breezes in adjacent squares”
except by writing one sentence for each square

First-order Logic

Propositional logic

Assumes that the world contains **facts**

First-order logic

Assumes that the world contains

- **Objects**
people, houses, numbers, theories, Donald Duck, colors, centuries, ...
- **Relations**
red, round, prime, multistoried, ...
brother of, bigger than, part of, has color, occurred after, owns, ...
- **Functions**
+, middle of, father of, one more than, beginning of, ...

Syntax of First-order Logic: Basic Elements

Symbols

Constants *KingJohn, 2, Koblenz, C, ...*

Predicates *Brother, >, =, ...*

Functions *Sqrt, LeftLegOf, ...*

Variables *x, y, a, b, ...*

Connectives $\wedge \vee \neg \Rightarrow \Leftrightarrow$

Quantifiers $\forall \exists$

Note

The **equality predicate** is always in the vocabulary
It is written in infix notation

Syntax of First-order Logic: Atomic Sentences

Atomic sentence

$predicate (term_1, \dots, term_n)$

or

$term_1 = term_2$

Term

$function (term_1, \dots, term_n)$

or

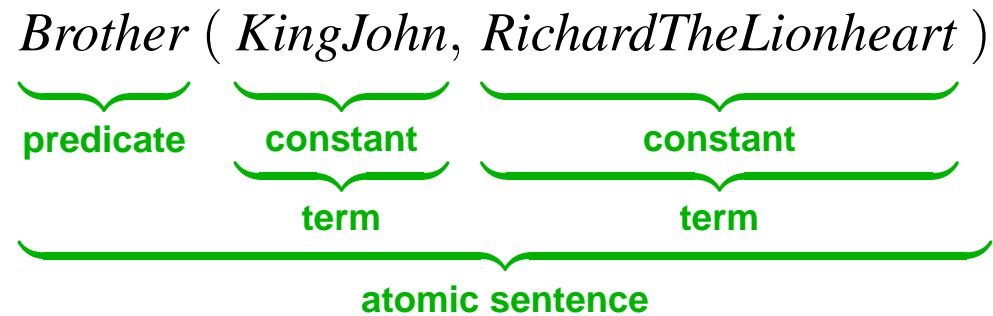
$constant$

or

$variable$

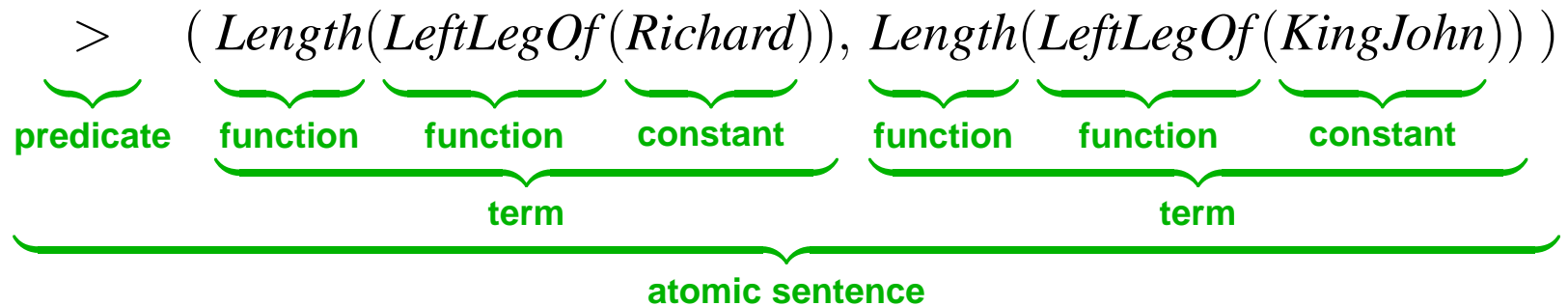
Syntax of First-order Logic: Atomic Sentences

Example



Syntax of First-order Logic: Atomic Sentences

Example



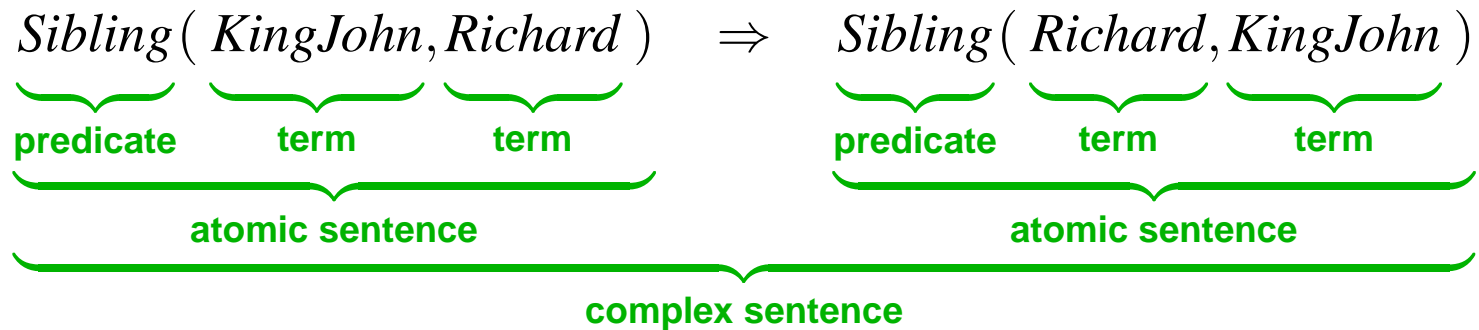
Syntax of First-order Logic: Complex Sentences

Built from atomic sentences using connectives

$$\neg S \quad S_1 \wedge S_2 \quad S_1 \vee S_2 \quad S_1 \Rightarrow S_2 \quad S_1 \Leftrightarrow S_2$$

(as in propositional logic)

Example



Semantics in First-order Logic

Models of first-order logic

Sentences are true or false with respect to models, which consist of

- a **domain** (also called universe)
- an **interpretation**

Domain

A non-empty (finite or infinite) set of arbitrary elements

Interpretation

Assigns to each

- constant symbol: a domain element
- predicate symbol: a relation on the domain (of appropriate arity)
- function symbol: a function on the domain (of appropriate arity)

Semantics in First-order Logic

Definition

An **atomic sentence**

$$predicate (term_1, \dots, term_n)$$

is true in a certain model (that consists of a domain and an interpretation)

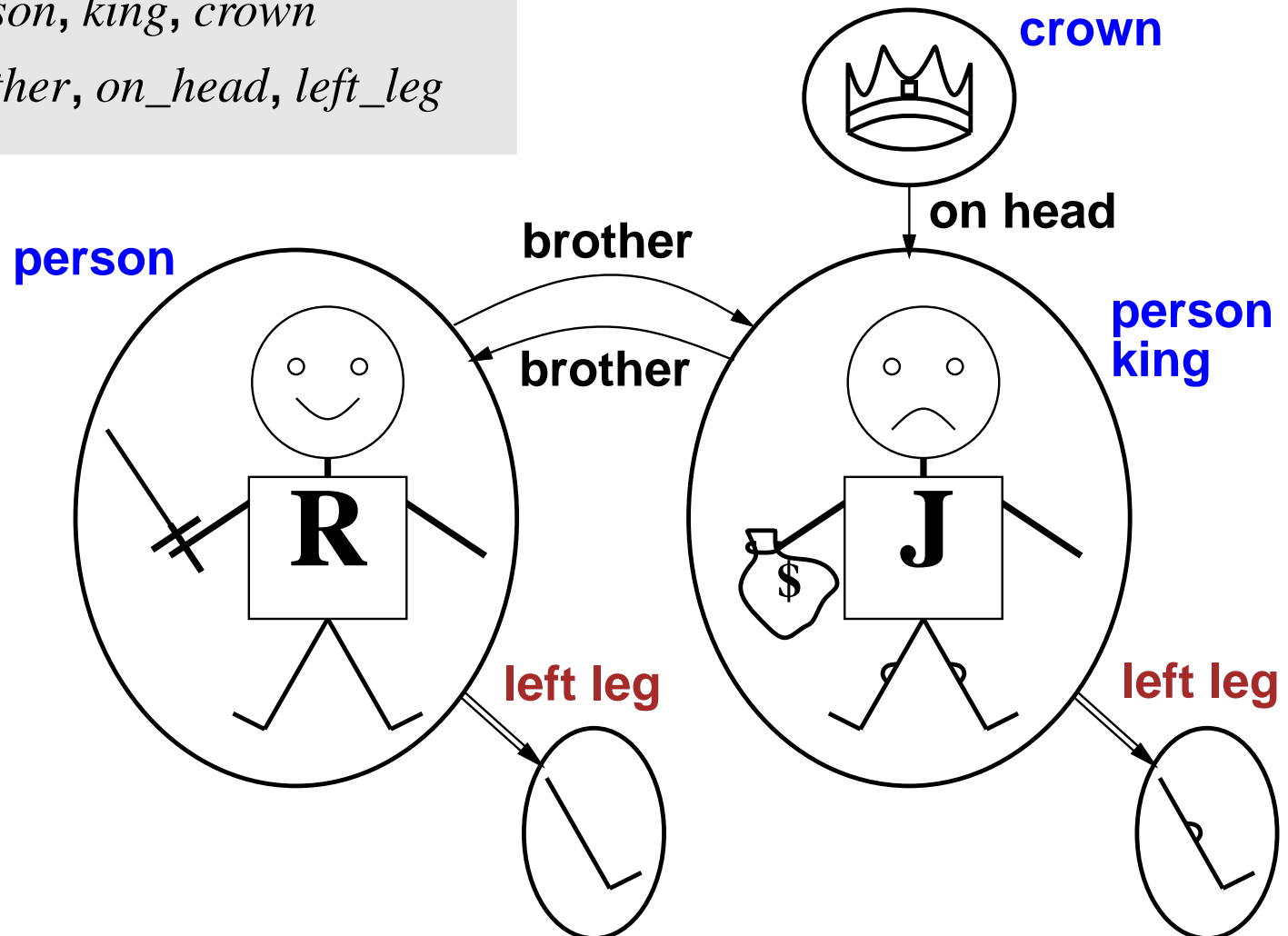
iff

the domain elements that are the interpretations of $term_1, \dots, term_n$
are in the relation that is the interpretation of *predicate*

The truth value of a **complex sentence** in a model
is computed from the truth-values of its atomic sub-sentences
in the same way as in propositional logic

Models for First-order Logic: Example

Constants: *KingJohn, Richard*
Predicates: *person, king, crown*
Functions: *brother, on_head, left_leg*



Universal Quantification: Syntax

Syntax

\forall *variables sentence*

Example

“Everyone studying in Koblenz is smart:

$$\forall \underbrace{x}_{\text{variables}} \underbrace{(StudiesAt(x, Koblenz) \Rightarrow Smart(x))}_{\text{sentence}}$$

Universal Quantification: Semantics

Semantics

$\forall xP$ is true in a model

iff

for all domain elements d in the model:

P is true in the model when x is interpreted by d

Intuition

$\forall xP$ is roughly equivalent to the conjunction of all instances of P

Example $\forall x \text{StudiesAt}(x, \text{Koblenz}) \Rightarrow \text{Smart}(x)$ equivalent to:

$\text{StudiesAt}(\text{KingJohn}, \text{Koblenz}) \Rightarrow \text{Smart}(\text{KingJohn})$

$\wedge \text{StudiesAt}(\text{Richard}, \text{Koblenz}) \Rightarrow \text{Smart}(\text{Richard})$

$\wedge \text{StudiesAt}(\text{Koblenz}, \text{Koblenz}) \Rightarrow \text{Smart}(\text{Koblenz})$

$\wedge \dots$

A Common Mistake to Avoid

Note

\Rightarrow is the main connective with \forall

Common mistake

Using \wedge as the main connective with \forall

Example

Correct: $\forall x (StudiesAt(x, Koblenz) \Rightarrow Smart(x))$

“Everyone who studies at Koblenz is smart”

Wrong: $\forall x (StudiesAt(x, Koblenz) \wedge Smart(x))$

“Everyone studies at Koblenz and is smart”, i.e.,

“Everyone studies at Koblenz and everyone is smart”

Existential Quantification: Syntax

Syntax

\exists *variables sentence*

Example

“Someone studying in Landau is smart:

$$\exists \underbrace{x}_{\text{variables}} \underbrace{(StudiesAt(x, Landau) \wedge Smart(x))}_{\text{sentence}}$$

Existential Quantification: Semantics

Semantics

$\exists xP$ is true in a model

iff

there is a domain element d in the model such that:

P is true in the model when x is interpreted by d

Intuition

$\exists xP$ is roughly equivalent to the disjunction of all instances of P

Example $\exists x \text{StudiesAt}(x, \text{Landau}) \wedge \text{Smart}(x)$ **equivalent to:**

$\text{StudiesAt}(\text{KingJohn}, \text{Landau}) \wedge \text{Smart}(\text{KingJohn})$

∨ $\text{StudiesAt}(\text{Richard}, \text{Landau}) \wedge \text{Smart}(\text{Richard})$

∨ $\text{StudiesAt}(\text{Landau}, \text{Landau}) \wedge \text{Smart}(\text{Landau})$

∨ ...

Another Common Mistake to Avoid

Note

\wedge is the main connective with \exists

Common mistake

Using \Rightarrow as the main connective with \exists

Example

Correct: $\exists x (StudiesAt(x, Landau) \wedge Smart(x))$

“There is someone who studies at Landau and is smart”

Wrong: $\exists x (StudiesAt(x, Landau) \Rightarrow Smart(x))$

“There is someone who, if he/she studies at Landau, is smart”

This is true if there is anyone not studying at Landau

Properties of Quantifiers

Quantifiers of same type commute

$\forall x \forall y$ **is the same as** $\forall y \forall x$

$\exists x \exists y$ **is the same as** $\exists y \exists x$

Properties of Quantifiers

Quantifiers of different type do NOT commute

$\exists x \forall y$ is **not** the same as $\forall y \exists x$

Example

$\exists x \forall y \text{Loves}(x, y)$

“There is a person who loves everyone in the world”

$\forall y \exists x \text{Loves}(x, y)$

“Everyone in the world is loved by at least one person”

(Both hopefully true but different)

Properties of Quantifiers

Quantifiers of different type do NOT commute

$\exists x \forall y$ is **not** the same as $\forall y \exists x$

Example

$\forall x \exists y \text{Mother}(x, y)$

“Everyone has a mother” **(correct)**

$\exists y \forall x \text{Mother}(x, y)$

“There is a person who is the mother of everyone” **(wrong)**

Properties of Quantifiers

Quantifier duality

$\forall x Likes(x, IceCream)$ **is the same as** $\neg \exists x \neg Likes(x, IceCream)$

$\exists x Likes(x, Broccoli)$ **is the same as** $\neg \forall x \neg Likes(x, Broccoli)$

Fun with Sentences

- **“Brothers are siblings”**

$$\forall x, y (Brother(x, y) \Rightarrow Sibling(x, y))$$

- **“Sibling” is symmetric**

$$\forall x, y (Sibling(x, y) \Leftrightarrow Sibling(y, x))$$

- **“One’s mother is one’s female parent”**

$$\forall x, y (Mother(x, y) \Leftrightarrow (Female(x) \wedge Parent(x, y)))$$

- **“A first cousin is a child of a parent’s sibling”**

$$\forall x, y (FirstCousin(x, y) \Leftrightarrow \exists p, ps (Parent(p, x) \wedge Sibling(ps, p) \wedge Parent(ps, y)))$$

Equality

Semantics

$term_1 = term_2$ is true under a given interpretation

if and only if

$term_1$ and $term_2$ have the same interpretation

Equality

Example

Definition of (full) sibling in terms of *Parent*

$$\begin{aligned} \forall x, y \text{ Sibling}(x, y) \Leftrightarrow & (\neg(x = y) \wedge \\ & \exists m, f (\neg(m = f) \wedge \\ & \text{Parent}(m, x) \wedge \text{Parent}(f, x) \wedge \\ & \text{Parent}(m, y) \wedge \text{Parent}(f, y))) \end{aligned}$$

Properties of First-order Logic

Important notions

- validity
- satisfiability
- unsatisfiability
- entailment

are defined for first-order logic in the same way as for propositional logic

Calculi

There are sound and complete calculi for first-order logic (e.g. resolution)

- Whenever $KB \vdash \alpha$, it is also true that $KB \models \alpha$
- Whenever $KB \models \alpha$, it is also true that $KB \vdash \alpha$

But these calculi **CANNOT decide** validity, entailment, etc.

Properties of First-order Logic

In propositional logic

Validity, satisfiability, unsatisfiability are **decidable**

In first-order logic

The set of valid, and the set of unsatisfiable formulas are **enumerable**

The set of satisfiable formulas is **NOT EVEN enumerable**